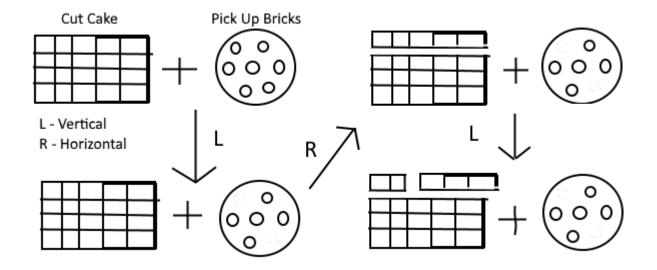
# $1 \ 2/15/17$

## 1.1 Sums of Position

• already used sums of component

Def'n:  $\delta, \alpha, \beta$  are positions in normal play games. Define  $\alpha + \beta$  to be a new position consisting of components  $\alpha$  and  $\beta$ . To move in  $\alpha + \beta$ , a player chooses a component to move in.

• A player moves from  $\alpha + \beta$  to either  $\alpha' + \beta$  or  $\alpha + \beta'$ .



#### 1.2 Determinate Sums

What is our strategy/ how do types behave under sums? Determine the type of the sum based on the type of the components.

Ex. Let  $\alpha$  be some position in some game of type R  $\alpha$  + P\cupB



What is R's move?

- Since the  $P \cup B$  game is type P, it doesn't change anything.
- Ignore  $P \cup B$  game until L makes a move. Then respond in that component.
- Normal play games last move wins, so next player plas in other comp.

Prop. If  $\beta$  is type P, then  $\alpha$  and  $\alpha + \beta$  are the same type.

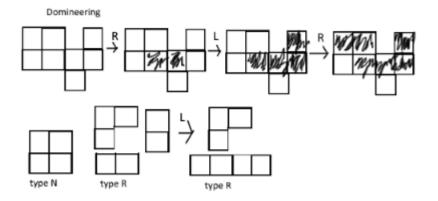
Prop. If  $\alpha$  and  $\beta$  are both type L(R) then  $\alpha + \beta$  is type L(R).

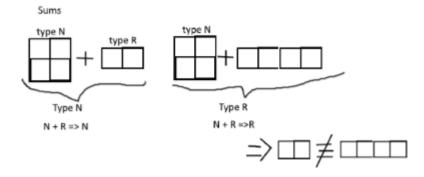
Т	L	R	Ν	Р
L	L	?	?	L
R	?	R	?	?
L R N P	?	?	?	Ν
Ρ	L	R	Ν	Ρ

### 1.3 Intermediate Sums

Domineering - cover some grid with dominoes which cover 2 spots at once.

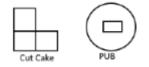
Assume R is horizontal, and L is vertical.





#### 1.4 Equivalence

Move beyond one game at a time



The next player to move will win either game.

Def'n: Two positions  $\alpha$  and  $\alpha'$  in (possibly different) normal play games are equivalent if for every position  $\beta$  in any normal play games, the two positions  $\alpha + \beta$  and  $\alpha' + \beta$  have the same type.

Equivalence Relations

Prop. If  $\alpha$ ,  $\beta$ ,  $\delta$  are positions in normal play games.

- 1.  $\alpha \equiv \alpha$  (reflexivity)
- 2.  $\alpha \equiv \beta \rightarrow \beta = \alpha$  (symmetry)
- 3.  $\alpha \equiv \beta$  and  $\beta \equiv \delta \rightarrow \alpha \equiv \delta$  (transitivity)

How does equivalence relate to type?

• If two positions are equivalent  $\rightarrow$  they have the same type

Prop. If  $\alpha \equiv \alpha'$ , then they have the same type.

• Let  $\beta$  be a position in a normal play game with no moves left.  $\alpha \leftrightarrow \alpha + \beta \leftrightarrow \alpha' + \beta \leftrightarrow \alpha'$ 

- Algebra w/ +,  $\equiv$
- Prop. If  $\alpha,\,\beta,\,\delta$  are positions in normal play games.
  - 1.  $\alpha + \beta \equiv \beta + \alpha$  (commutably)
  - 2.  $(\alpha + \beta) + \delta \equiv \alpha + (\beta + \delta)$  (associatably)

Lemma. Given position  $\alpha$ ,  $\beta$ , $\delta$  in normal play games

- 1. If  $\alpha \equiv \alpha'$ , then  $\alpha + \beta \equiv \alpha' + \beta$
- 2. If  $\alpha_i \equiv \alpha_i$  for  $1 \leq i \leq n$ , then  $\alpha_i + \ldots + \alpha_n \equiv \alpha_i + \alpha_n$
- 3. If  $\alpha_i \equiv \alpha_i$  for  $1 \leq i \leq m$  and  $\beta_j \equiv \beta_j$ ' for  $1 \leq j \leq n$ , then  $\{\alpha_n, \alpha_m \beta_1, ..., \beta_n\} \equiv \{\alpha_1, ..., \alpha_n' \beta_1, ..., \beta_n\}$

Type P is equivalent to zero under normal addition. Lemma. If  $\beta$  is type P, then  $\alpha + \beta = \alpha$ 

Prop. If  $\alpha$  and  $\alpha'$  are type P, then  $\alpha \equiv \alpha' \rightarrow \alpha + \delta \equiv \delta \equiv \alpha' + \delta$ 

Lemma. If  $\alpha + \beta$  and  $\alpha' + \beta$  are both type P, then  $\alpha \equiv \alpha'$ 

### 1.5

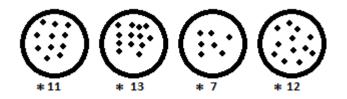
Cor. Every position in an impartial game is one of the following:

- Type N The next player to play has a winning strategy.
- Type P The previous (or second) player to play has a winning strategy.

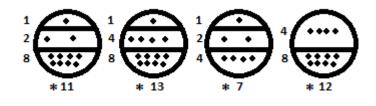
Cor. A position in an impartial game is:

- Type N If  $\exists$  a move to a position of type P.
- Type P If there is no move to a position of type P.

For convenience, # of stones in each pile



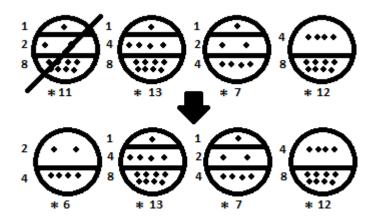
Def'n: Given a Nim position  $*\alpha_1 + *\alpha_2 + ... + *\alpha_k$ , it is balanced if, for every power of 2, the total # of subpiles of that size is even.



 $2^0+2^1+2^2+\ldots+2^n<2^{n+1}$ 

Procedure - Balance an unbalanced position:

- Let  $*\alpha_1 + *\alpha_2 + \dots + *\alpha_k$  be an unbalanced game of Nim.
- Suppose  $2^n$  is the largest power of 2 which there are an odd # of subpiles Curretly all subpiles of  $2^i, j > m$  are balanced for every j; m, if there are an odd # of subpiles of  $S^j$ , excluding this pile, leave  $s^i$  stones. Since at least  $2^m$  stones could be removed, and  $2^0 + 2^i + \ldots + 2^{n-1} < 2^m$  stones should be left, this is a legal move.



Prop. Every balanced Nim position is type P and every unbalanced Nim position is type N.

#### 1.6 Nimbers

Its helpful to think of Nim positions like #'s

- we call a stack  $*\alpha$  a nimber
- $\bullet$  Based on prop and that P doesn't additively change anything a type P position \*0
- $\alpha \equiv \gamma$
- every balanced position  $\gamma \equiv *0$

Imagine the  $2^{nd}$  player could add a stack to the game before it starts Original  $*a_1+\ldots+*a_L$ 

- $\bullet$  add \*b
- to a balanced position of type P to determine odd # of subpiles of  $2^j$

 $*a_1 + \ldots + *a_l + *b = *0$