

Problem 1: Let $\text{RELPRIME} = \{\langle x, y \rangle \mid x \text{ and } y \text{ are positive integers that are relatively prime}\}$. Given the following algorithm to test if two positive integers are relatively prime, let n be the maximum number of decimal digits in x and y . Analyze the running time of this algorithm, using O -notation. Explain the details and give your reasoning for each step.

On input $\langle x, y \rangle$ where x and y are positive integers.

1. Repeat until $y = 0$:
2. Assign $x \leftarrow x \bmod y$.
3. Swap x and y .
4. Output x . If the result is 1, *accept*; otherwise *reject*.

Problem 2: Prove the following problems are in **P** by giving an algorithm and showing their run-times.

- (a) Path = $\{\langle G, s, t \rangle : \text{Graph } G \text{ has a path from vertex } s \text{ to vertex } t\}$.
- (b) No-Path = $\{\langle G, s, t \rangle : \text{Graph } G \text{ does not have a path from vertex } s \text{ to vertex } t\}$.
- (c) Euler-Path = $\{\langle G \rangle : \text{Graph } G \text{ has an Euler path}\}$.
- (d) Maximum-Matching = $\{\langle G \rangle : \text{A maximal matching for Graph } G\}$.
- (e) 1-pos-3-SAT = $\{\langle \phi \rangle : \text{Boolean formula in 3-CNF form has a satisfying assignment and each clause has a variable that is not negated}\}$.

Problem 3: Prove the following problems are in **NP** by giving a nondeterministic TM, and also by giving a deterministic verifier.

- (a) Harnpath = $\{\langle G, s, t \rangle : \text{Graph } G \text{ has a Hamiltonian path from vertex } s \text{ to vertex } t\}$.
- (b) Vertex-Cover = $\{\langle G, k \rangle : \text{Graph } G \text{ has a vertex cover of size } k\}$.
- (c) Traveling-Salesman-Problem, TSP = $\{\langle G, k \rangle : \text{Weighted graph } G \text{ has a tour of length } k\}$.
- (d) Bin-Packing, BP = $\{\langle S, c, k \rangle : S \text{ is a set of numbers that can be put into } k \text{ bins of size } c.\}$
- (e) Minimum-Maximal-Matching, MMM = $\{\langle G, k \rangle : \text{Graph } G \text{ has a maximal matching of size } k\}$.
- (f) 1-in-3-SAT = $\{\langle \phi \rangle : \text{Boolean formula in 3-CNF form has a satisfying assignment and each clause only has one variable that is true}\}$.

Problem 4: Given an algorithm that can solve a decision problem in $O(t(n))$ time (an oracle), show that we can also solve the optimization version with only a polynomial difference.