

Problem 1: Prove the following problems are in **L** by giving an algorithm and showing their run-times.

- (a) Path = $\{ \langle G, s, t \rangle : \text{Graph } G \text{ has a path from vertex } s \text{ to vertex } t \}$.
- (b) Majority = $\{ \langle w \rangle : \text{the number of 1's in the string is greater than the number of 0's} \}$.
- (c) Proper-parentheses = $\{ \langle w \rangle : w \text{ is a string of properly nested parentheses} \}$.

Problem 2: Show that 2SAT is in **NL** by giving an algorithm to solve it using nondeterministic log-space.

Problem 3: Prove the following problems are **NP**-hard by giving a polynomial-time reduction from another NP-hard problem.

- (a) HamCycle = $\{ \langle G \rangle : \text{Graph } G \text{ has a Hamiltonian cycle} \}$ (reduce from *Undirected Hamiltonian Path*)
- (b) 2-m-SAT = $\{ \langle \phi, k \rangle : \phi \text{ is satisfiable with only } k \text{ variables set to true} \}$ (reduce from *Vertex Cover*)
- (c) Traveling-Salesman-Problem, TSP = $\{ \langle G, k \rangle : \text{Weighted graph } G \text{ has a tour of length } k \}$ (reduce from *Undirected Hamiltonian Path*)
- (d) Bin-Packing, BP = $\{ \langle S, c, k \rangle : S \text{ is a set of numbers that can be put into } k \text{ bins of size } c \}$ (reduce from the *Partition Problem*)
- (e) Hitting-Set, HS = $\{ \langle S, k \rangle : S \text{ is a set of sets of numbers, and there is a set of size } k \text{ that hits every set in } S \}$ (reduce from *Vertex Cover*)
- (f) Longest-Path = $\{ \langle G, a, b, k \rangle \mid G \text{ contains a simple path of length at least } k \text{ from } a \text{ to } b \}$ (reduce from *Undirected Hamiltonian Path*)

Problem 4: Explain why problems like the Partition problem are only considered *weakly* **NP**-hard. What is pseudopolynomial time?